

## Information Theory

### Exercise Sheet 5

**Exercise 1** (About the binomial).

1. Show with the method of types that,

$$\frac{1}{(n+1)^2} 2^{nh(k/n)} \leq \binom{n}{k} \leq 2^{nh(k/n)}$$

2. Our main goal now is to sharpen the lower-bound by showing that

$$\frac{1}{n+1} 2^{nh(k/n)} \leq \binom{n}{k} \leq 2^{nh(k/n)}$$

*du si d'u s'parameters n, p is up*

**Hint:** first show that the most likely event of a binomial distribution

**Exercise 2** (Relation between  $D_{\text{KL}}(P||Q)$  and chi-square). Show that the  $\chi^2$  statistic

$$\chi^2 = \sum_x \frac{(P(x) - Q(x))^2}{Q(x)}$$

is twice the first term in the Taylor series expansion of  $D_{\text{KL}}(P||Q)$  about  $Q$ .

**Hint:** Write  $P/Q = 1 + \epsilon$  and expand the  $\log_2$

**Exercise 3** (Error exponent for universal codes).

1. Show that the universal source code of lecture with rate  $R$  achieves a probability of error  $P_e^{(n)} \stackrel{\text{poly}}{=} 2^{-nD_{\text{KL}}(P^*||Q)}$ , where  $Q$  is the true distribution and  $P^*$  achieves  $\min D_{\text{KL}}(P||Q)$  over all the  $P$ 's such that  $H(P) \geq R$
2. Find  $P^*$  in terms of  $Q$  and  $R$
3. Now let  $\mathbf{X} \in \{0, 1\}$ . Find the region of source probabilities  $Q(x)$  for which rate  $R$  is sufficient for the universal source code to achieve  $P_e^{(n)} \rightarrow 0$ .

**Exercise 4.** Consider a joint distribution  $Q(x, y)$  with marginals  $Q(x)$  and  $Q(y)$ . Let  $E$  be the set of types that look jointly typical with respect to  $Q$ :

$$E \stackrel{\text{def}}{=} \left\{ P(x, y) \in \mathcal{P} : \begin{aligned} & \left| - \sum_{x,y} P(x, y) \log_2 Q(x) - H(\mathbf{X}) \right| = 0, \\ & \left| - \sum_{x,y} P(x, y) \log_2 Q(y) - H(\mathbf{Y}) \right| = 0, \\ & \left| - \sum_{x,y} P(x, y) \log_2 Q(x, y) - H(\mathbf{X}, \mathbf{Y}) \right| = 0 \end{aligned} \right\}$$

1. Let  $Q_0(x, y)$  be another distribution on  $\mathcal{X} \times \mathcal{Y}$ . Argue that the distribution  $P^*$  in  $E$  that is the closest to  $Q_0$  is of the form

$$P^*(x, y) = Q_0(x, y) 2^{\lambda_0 + \lambda_1 \log_2 Q(x) + \lambda_2 \log_2 Q(y) + \lambda_3 \log_2 Q(x, y)} \quad (1)$$

where  $\lambda_0, \lambda_1, \lambda_2$  and  $\lambda_3$  are chosen to satisfy the constraints. Argue that this distribution is unique.

2. Now let  $Q_0(x, y) = Q(x)Q(y)$ . Verify that  $Q(x, y)$  is of the form (1) and that it belongs to  $E$ . Deduce that  $P^* = Q$ .

**Exercise 5 (Counting).** Let  $\mathcal{X} = \{1, 2, \dots, m\}$ . Show that the number of sequences  $\mathbf{x} \in \mathcal{X}^n$  satisfying  $1/n \sum_{i=1}^n g(x_i) \geq \alpha$  is approximately equal to  $2^{nH^*}$ , to first order in the exponent, for  $n$  sufficiently large, where

$$H^* = \max_{P: \sum_{i=1}^m P(i)g(i) \geq \alpha} H(P)$$

**Exercise 6 (Running difference).** Let  $\mathbf{X}_1, \dots, \mathbf{X}_n$  be i.i.d. according to  $Q_1$ , and  $\mathbf{Y}_1, \dots, \mathbf{Y}_n$  be i.i.d. according to  $Q_2$ . Let  $\mathbf{X}^n$  and  $\mathbf{Y}^n$  be independent. Find an expression for

$$\mathbb{P} \left( \sum_{i=1}^n \mathbf{X}_i - \mathbf{Y}_i \geq nt \right)$$

good to first order in the exponent.

**Exercise 7** (Type constraint).

1. Find constraints on the type  $P_{\mathbf{X}^n}^{\text{emp}}$  such that the sample variance verifies

$$\overline{\mathbf{X}_n^2} - (\overline{\mathbf{X}_n})^2 \leq \alpha$$

where

$$\overline{\mathbf{X}_n^2} \stackrel{\text{def}}{=} \frac{1}{n} \sum_{i=1}^n \mathbf{X}_i^2 \quad \text{and} \quad \overline{\mathbf{X}_n} \stackrel{\text{def}}{=} \frac{1}{n} \sum_{i=1}^n \mathbf{X}_i$$

2. Find the exponent in the probability  $Q^n \left( \overline{\mathbf{X}_n^2} - (\overline{\mathbf{X}_n})^2 \leq \alpha \right)$ .